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**APPARATUS AND METHOD FOR ENCRYPTING AND DECRYPTING DATA
WITH INCREMENTAL DATA VALIDATION**

BACKGROUND OF THE INVENTION

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1. Technical Field:

The present invention is directed to an improved computing device. More specifically, the present invention is directed to an apparatus and method for encrypting and
10 decrypting data with incremental data validation.

2. Description of Related Art:

Internet Protocols which use cryptography are prone to Denial of Service (DOS) attacks because cryptography
15 requires a large amount of processor time. A DOS attack is an assault on a network that floods it with so many additional requests that regular traffic is either slowed or completely interrupted. The regular traffic is slowed or completely interrupted because the victim computer systems
20 must expend resources to decrypt the data in these numerous requests only to find that the requests are not authentic. Thus, resources that could be used to handle regular traffic is instead tied up with handling unauthentic requests sent as part of a DOS attack.

25 In order to avoid such attacks, messages and packets which are encrypted may have a digital digest attached to them for authentication purposes. A digital digest is a mechanism used to uniquely identify the contents of the message or packet. A digital digest may be a checksum or
30 the like, for example.

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Figure 1 is a diagram illustrating a known mechanism for encrypting data. As shown in **Figure 1**, clear text data **110** is initially received. The data is encrypted to product encrypted data **120**. Encrypted data is read byte by byte to
5 create a unique digital digest **130** for the encrypted data. The digital digest is encrypted and appended to the encrypted data to thereby produce and encrypted message or packet **140**. The encrypted message or packet **140** may then be transmitted to a receiving device.

10 At the receiving device, in order to process the data, the message or packet **140** must first be authenticated and decrypted before the processor is able to process the encrypted data. In order to authenticate the message or packet **140**, all of the encrypted data **120** in the message or
15 packet **140** must first be read to calculate a corresponding digital digest. The digital digest **130** appended to the encrypted data **120** is then decrypted and compared to the digital digest calculated based on the encrypted data in the received data message or packet **140**.

20 If the two digital digests match, the data message or packet **140** is authentic. If the data message or packet **140** is authentic, then the encrypted data **120** may be decrypted and processed. Otherwise, if the data message or packet **140** is not authentic, the data message or packet **140** is
25 discarded. Thus, with the prior art mechanisms, all of the encrypted data in the data message or packet **140** must be read twice in order to authenticate and decrypt the data message or packet **140**.

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Therefore, it would be beneficial to have an apparatus and method by which data messages or packets may be authenticated and decrypted using a single pass on the encrypted data. Moreover, it would be beneficial to have an
5 apparatus and method for incrementally authenticating a data message or packet based on a digital digest so that processing of non-authentic data messages or packets is halted at an earliest possible time to thereby free resources that may be used in authenticating and decrypting
10 authentic data messages or packets.

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SUMMARY OF THE INVENTION

5 The present invention provides an apparatus and method for encrypting and decrypting data with incremental data validation. With the mechanism of the present invention, data is encrypted and a digital digest is generated in chunks. That is, the digital digest is comprised of a plurality of intermediate digital digest chunks, each of which can be used to validate a portion of the associated
10 encrypted data. During decryption, a portion of the encrypted data is read and decrypted at approximately the same time that a digital digest is calculated for that portion of the encrypted data.

15 The calculated partial digital digest may then be compared to an intermediate digital digest associated with the portion of the encrypted data, and which is appended to the encrypted data. If the two digital digests match, decryption of the encrypted data may proceed to the next portion of the encrypted data. If the two digital digests
20 do not match, decryption is halted and the data message or packet is discarded without having decrypted the entire data message or packet.

25 In this way, resources may be freed from processing non-authentic data messages or packets so that they may be used in processing authentic data messages. Thus, the susceptibility of the present invention to denial of service attacks is noticeably reduced in comparison with the prior art.

BRIEF DESCRIPTION OF THE DRAWINGS

The novel features believed characteristic of the invention are set forth in the appended claims. The invention itself, however, as well as a preferred mode of use, further objectives and advantages thereof, will best be understood by reference to the following detailed description of an illustrative embodiment when read in conjunction with the accompanying drawings, wherein:

Figure 1 is an exemplary diagram of a prior art method of encrypting/decrypting data using a digital digest;

Figure 2 is an exemplary diagram illustrating a distributed data processing system in accordance with the present invention;

Figure 3 is an exemplary diagram illustrating a server data processing device in accordance with the present invention;

Figure 4 is an exemplary diagram illustrating a client data processing device in accordance with the present invention;

Figure 5 is a diagram illustrating an encryption operation according to the present invention;

Figure 6 is a diagram illustrating a decryption operation according to the present invention;

Figure 7 is a flowchart outlining an exemplary operation for encrypting data according to the present invention; and

Figure 8 is a flowchart outlining an exemplary operation for decrypting data according to the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

With reference now to the figures, **Figure 2** depicts a pictorial representation of a network of data processing systems in which the present invention may be implemented. Network data processing system **200** is a network of computers in which the present invention may be implemented. Network data processing system **200** contains a network **202**, which is the medium used to provide communications links between various devices and computers connected together within network data processing system **200**. Network **202** may include connections, such as wire, wireless communication links, or fiber optic cables.

In the depicted example, server **204** is connected to network **202** along with storage unit **206**. In addition, clients **208**, **210**, and **212** are connected to network **202**. These clients **208**, **210**, and **212** may be, for example, personal computers or network computers. In the depicted example, server **204** provides data, such as boot files, operating system images, and applications to clients **208-212**. Clients **208**, **210**, and **212** are clients to server **204**. Network data processing system **200** may include additional servers, clients, and other devices not shown.

In the depicted example, network data processing system **200** is the Internet with network **202** representing a worldwide collection of networks and gateways that use the TCP/IP suite of protocols to communicate with one another. At the heart of the Internet is a backbone of high-speed data communication lines between major nodes or host computers, consisting of thousands of commercial, government, educational and other computer systems that

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route data and messages. Of course, network data processing system **200** also may be implemented as a number of different types of networks, such as for example, an intranet, a local area network (LAN), or a wide area network (WAN). **Figure 2** is intended as an example, and not as an architectural limitation for the present invention.

Referring to **Figure 3**, a block diagram of a data processing system that may be implemented as a server, such as server **204** in **Figure 2**, is depicted in accordance with a preferred embodiment of the present invention. Data processing system **300** may be a symmetric multiprocessor (SMP) system including a plurality of processors **302** and **304** connected to system bus **306**. Alternatively, a single processor system may be employed. Also connected to system bus **306** is memory controller/cache **308**, which provides an interface to local memory **309**. I/O bus bridge **310** is connected to system bus **306** and provides an interface to I/O bus **312**. Memory controller/cache **308** and I/O bus bridge **310** may be integrated as depicted.

Peripheral component interconnect (PCI) bus bridge **314** connected to I/O bus **312** provides an interface to PCI local bus **316**. A number of modems may be connected to PCI local bus **316**. Typical PCI bus implementations will support four PCI expansion slots or add-in connectors. Communications links to network computers **208-212** in **Figure 2** may be provided through modem **318** and network adapter **320** connected to PCI local bus **316** through add-in boards.

Additional PCI bus bridges **322** and **324** provide interfaces for additional PCI local buses **326** and **328**, from which additional modems or network adapters may be

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supported. In this manner, data processing system **300** allows connections to multiple network computers. A memory-mapped graphics adapter **330** and hard disk **332** may also be connected to I/O bus **312** as depicted, either
5 directly or indirectly.

Those of ordinary skill in the art will appreciate that the hardware depicted in **Figure 3** may vary. For example, other peripheral devices, such as optical disk drives and the like, also may be used in addition to or in place of the
10 hardware depicted. The depicted example is not meant to imply architectural limitations with respect to the present invention.

The data processing system depicted in **Figure 3** may be, for example, an IBM e-Server pSeries system, a product of
15 International Business Machines Corporation in Armonk, New York, running the Advanced Interactive Executive (AIX) operating system or LINUX operating system.

With reference now to **Figure 4**, a block diagram illustrating a data processing system is depicted in which
20 the present invention may be implemented. Data processing system **400** is an example of a client computer. Data processing system **400** employs a peripheral component interconnect (PCI) local bus architecture. Although the depicted example employs a PCI bus, other bus architectures
25 such as Accelerated Graphics Port (AGP) and Industry Standard Architecture (ISA) may be used. Processor **402** and main memory **404** are connected to PCI local bus **406** through PCI bridge **408**. PCI bridge **408** also may include an integrated memory controller and cache memory for processor
30 **402**. Additional connections to PCI local bus **406** may be made through direct component interconnection or through

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add-in boards. In the depicted example, local area network (LAN) adapter **410**, SCSI host bus adapter **412**, and expansion bus interface **414** are connected to PCI local bus **406** by direct component connection. In contrast, audio adapter **416**, graphics adapter **418**, and audio/video adapter **419** are connected to PCI local bus **406** by add-in boards inserted into expansion slots. Expansion bus interface **414** provides a connection for a keyboard and mouse adapter **420**, modem **422**, and additional memory **424**. Small computer system interface (SCSI) host bus adapter **412** provides a connection for hard disk drive **426**, tape drive **428**, and CD-ROM drive **430**. Typical PCI local bus implementations will support three or four PCI expansion slots or add-in connectors.

An operating system runs on processor **402** and is used to coordinate and provide control of various components within data processing system **400** in **Figure 4**. The operating system may be a commercially available operating system, such as Windows 2000, which is available from Microsoft Corporation. An object oriented programming system such as Java may run in conjunction with the operating system and provide calls to the operating system from Java programs or applications executing on data processing system **400**. "Java" is a trademark of Sun Microsystems, Inc. Instructions for the operating system, the object-oriented operating system, and applications or programs are located on storage devices, such as hard disk drive **426**, and may be loaded into main memory **404** for execution by processor **402**.

Those of ordinary skill in the art will appreciate that the hardware in **Figure 4** may vary depending on the implementation. Other internal hardware or peripheral

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devices, such as flash ROM (or equivalent nonvolatile memory) or optical disk drives and the like, may be used in addition to or in place of the hardware depicted in **Figure 4**. Also, the processes of the present invention may be applied to a multiprocessor data processing system.

As another example, data processing system **400** may be a stand-alone system configured to be bootable without relying on some type of network communication interface, whether or not data processing system **400** comprises some type of

network communication interface. As a further example, data processing system **400** may be a Personal Digital Assistant (PDA) device, which is configured with ROM and/or flash ROM in order to provide nonvolatile memory for storing operating system files and/or user-generated data.

The depicted example in **Figure 4** and above-described examples are not meant to imply architectural limitations. For example, data processing system **400** also may be a notebook computer or hand held computer in addition to taking the form of a PDA. Data processing system **400** also may be a kiosk or a Web appliance.

Figure 5 is an exemplary diagram illustrating a data encryption operation according to the present invention. The operation shown in **Figure 5** may be implemented as hardware, software, or a combination of hardware and software. For example, in a preferred embodiment, the present invention is implemented as software instructions executed by a processor on data stored in a memory, storage device, or buffer. For example, the present invention may be implemented as computer program instructions executed by one or more of the processors **302**, **304** and **402** on data stored in a memory, storage device or buffer, such as local memory **309**, hard

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disk **332**, main memory **404**, disk **426**, tape **428**, CD-ROM **430**,
memory **424**, or the like. Alternatively, the present
invention may be implemented using data obtained via a
communications interface such as modem **318**, network adapter
5 **320**, LAN adapter **410**, or modem **422**. Other embodiments of
the present invention may obtain data for use with the
present invention via other mechanisms without departing
from the spirit and scope of the present invention.

As shown in **Figure 5**, clear data **510** is read in chunks
10 and encrypted as a plurality of encrypted data portions
531-535. The encrypted data portions **531-535** correspond to
chunks of data and may be of any desirable size. In an
exemplary embodiment, the encrypted data portions **531-535**
correspond to 64 byte data chunks of the clear data **510**. In
15 an exemplary embodiment, the data is read and stored in a
buffer (not shown) which then outputs the data to a
processor in chunks of a predetermined size. As the chunks
of data are output from the buffer, the present invention is
implemented on the data chunks.

20 For each of the encrypted data portions **531-535**, a
digital digest is generated. The generation of a digital
digest from encrypted data is generally known in the art and
thus, a detailed explanation of the procedures for
generating a digital digest will not be provided herein.

25 The digital digests of the present invention, however,
differ from known digital digest generation mechanism in
that a digital digest is generated for one or more
intermediate portions of the encrypted data. In this way, a
plurality of intermediate digital digests are generated.

30 Each of the plurality of intermediate digital digests
are encrypted to thereby generate intermediate encrypted

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digital digests **541-545** which are appended to the end of the encrypted data message or packet **540**. Thus, the data message or packet **540** is comprised of a plurality of encrypted data portions **531-535** and corresponding

5 intermediate encrypted digital digests **541-545**.

Figure 6 is an exemplary diagram illustrating an operation for reading, authenticating, and decrypting the encrypted data message or packet **540** according to the present invention. As with the operation shown in **Figure 5**,

10 the operation shown in **Figure 6** may be implemented as software, hardware or a combination of software and hardware, depending on the particular embodiment.

As shown in **Figure 6**, the operation first reads a first encrypted data portion **610** and calculates a digital digest
15 **620** from the first encrypted data portion **610**. The operation then reads and decrypts an intermediate encrypted digital digest **541**, from the end of the data message or packet **540**, that corresponds to the first encrypted data portion **610**. The decrypted intermediate digital digest **630**
20 is then compared to the calculated digital digest **620**. If the two digital digests do not match, the data is not authentic or is otherwise corrupted and the data message or packet **540** is discarded.

If the two digital digests do match, the encrypted data
25 portion **610** is decrypted and the next encrypted data portion **640** is read from the data message or packet **540**. The process then continues in the same manner. At any time during the process, if any one of the digital digest comparisons results in a non-match, the data message or
30 packet **540** is discarded.

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Thus, the present invention provides a mechanism in which only a single pass through the encrypted data is necessary to both authenticate and decrypt the data. The present invention uses an incremental approach to

5 authenticate portions of the encrypted data and decrypt the data. If any one of the authentication procedures results in an indication that the data may be unauthentic or corrupted, the entire data message or packet is discarded. In this way, unauthentic or corrupted data is identified at
10 an earliest possible time during the authentication and decryption process. Therefore, resources are freed at an earlier time so that they may be used to authenticate and decrypt authentic and/or uncorrupted data.

Figure 7 is a flowchart outlining an exemplary
15 operation of the present invention when encrypting a data message or packet. As shown in **Figure 7**, the operation starts with reading the next data chunk of the data message or packet (step **710**). If this is the first time through the operation, the next data chunk is the first data chunk in
20 the data message or packet. The data chunk is then encrypted (step **720**) and an intermediate digital digest is generated for the encrypted data chunk (step **730**). This intermediate digital digest is preferably stored in memory until all data chunks of the data message or packet are
25 encrypted and the data message or packet is ready for transmission.

A determination is then made as to whether the data chunk is the last data chunk in the data message or packet (step **740**). If the data chunk is not the last data chunk in
30 the data message or packet, the operation returns to step **710** and performs steps **710-730** on the next data chunk in the

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data message or packet. If the data chunk is the last data chunk in the data message or packet, the intermediate digital digests are appended to the encrypted data (step 750) and the operation ends. The data message or packet is then ready for storage or transmission.

Figure 8 is a flowchart outlining an exemplary operation of the present invention when decrypting a data message or packet. As shown in **Figure 8**, the operation starts with reading the next portion of the encrypted data in the data message or packet (step 810). If this is the first time the operation is executed, the next portion of the encrypted data is a first portion of the encrypted data.

A digital digest is then calculated for the portion of the encrypted data (step 820). An appended intermediate digital digest corresponding to the portion of encrypted data is then decrypted (step 830) and compared to the calculated digital digest (step 840). A determination is then made as to whether the data is authentic based on the comparison (step 850).

If the data is not authentic, the entire data message or packet is discarded (step 880). If the data is authentic, the portion of encrypted data is decrypted and processing of the data message or packet is continued with the next portion of encrypted data in the data message or packet (step 860). A determination is made as to whether the portion is the last data portion in the data message or packet (step 870). If not, the operation returns to step 810. Otherwise, if the data portion is the last data portion in the data message or packet, the operation terminates.

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While the above embodiments of the present invention have been described in terms of a one-to-one correspondence between data chunks and intermediate digital digests, such a convention is used only for simplicity of illustration of the present invention. The present invention is not limited to such embodiments. Rather, the size of the data chunks and the size of data used to generate the digital digests may be different without departing from the spirit and scope of the present invention.

Furthermore, while the above embodiments have been described in terms of intermediate digital digests that correspond to separate portions of encrypted data in the data message or packet, the present invention is not limited to such embodiments. Rather, as an alternative embodiment, the portions of encrypted data may be built up in increments of chunks of data and the corresponding digital digests may likewise be built up. In other words, assume a data message is comprised of a first, second and third data chunk. The first portion of encrypted data would correspond to an encrypted first data chunk. The second portion of the encrypted data would correspond to an encrypted combination of the first and second data chunks. The third portion of the encrypted data would correspond to an encrypted combination of the first, second and third data chunks.

As a result, the intermediate digital digests would include a first intermediate digital digest calculated from the encrypted first data chunk. The second intermediate digital digest would be calculated from a combination of the encrypted first data chunk and an encrypted second data chunk. The third intermediate digital digest would be calculated from a combination of then encrypted first, second and third data chunks. Other mechanisms for setting

forth the data portions and the intermediate digital digests may be used without departing from the spirit and scope of the present invention.

Thus, the present invention provides a mechanism in which a data message or packet may be authenticated and decrypted with a single pass on the encrypted data. The present invention avoids the problems of the prior art by reducing the amount of operations necessary to perform authentication and decryption. Since the present invention is capable of identifying unauthentic data or corrupted data prior to decrypting the entire data message or packet, the present invention is less susceptible to denial of service attacks.

It is important to note that while the present invention has been described in the context of a fully functioning data processing system, those of ordinary skill in the art will appreciate that the processes of the present invention are capable of being distributed in the form of a computer readable medium of instructions and a variety of forms and that the present invention applies equally regardless of the particular type of signal bearing media actually used to carry out the distribution. Examples of computer readable media include recordable-type media such as floppy disc, a hard disk drive, a RAM, and CD-ROMs and transmission-type media such as digital and analog communications links.

The description of the present invention has been presented for purposes of illustration and description, but is not intended to be exhaustive or limited to the invention in the form disclosed. Many modifications and variations 30 will be apparent to those of ordinary skill in the art. The embodiment was chosen and described in order to best explain

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the principles of the invention, the practical application,
and to enable others of ordinary skill in the art to
understand the invention for various embodiments with
various modifications as are suited to the particular use
5 contemplated.

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